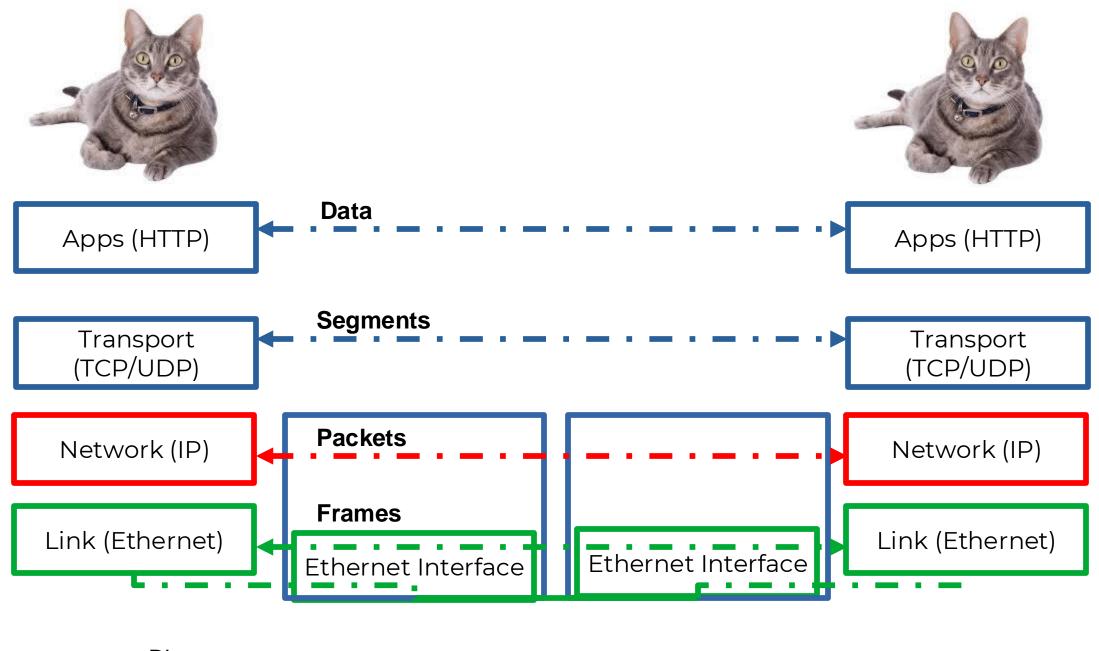
CSC2710 – Intro to System and Networking

Instructor: Susmit Shannigrahi

INTERNET PROTOCOL (IP) sshannigrahi@tntech.edu

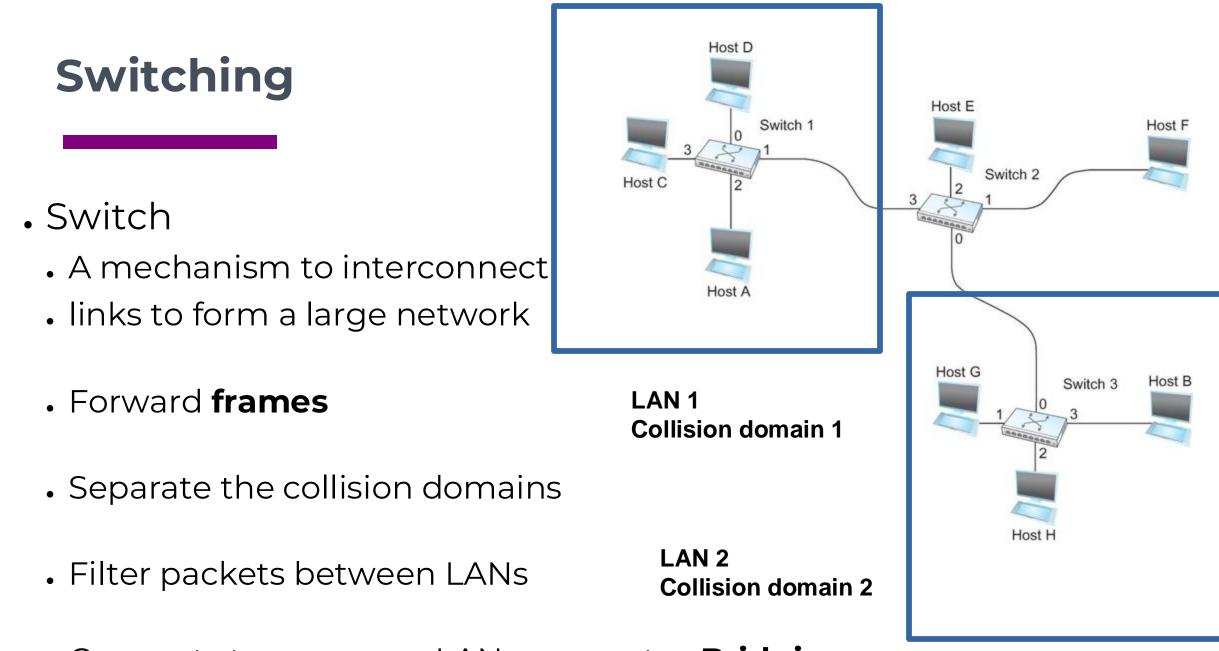




Bits (1010001)

So far...

- We are forwarding packets between different LANs
- Spanning tree algorithm for preventing loops

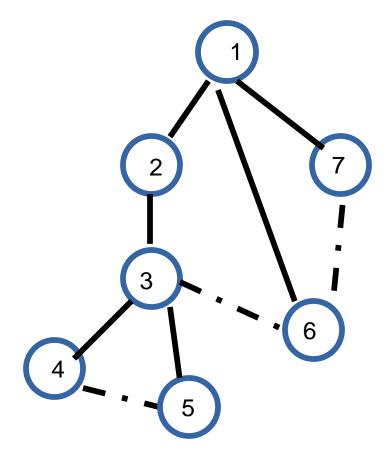


Connects two or more LAN segments - Bridging

How do we create a spanning tree?

. Message (Y, d, X) - (to, distance, from)

- 4 thinks it's the root
- Sends (4, 0, 4) to 3 and 5
- Receives (3,0,3) from 3
 - Sets it to as the root since 3 < 4
- Receives (3,1,5) from 5
 - Sees that this is a longer path to 3
 - 2 hops vs direct path (1 hop)
 - Removes 4-5 link from the tree
- . Does not scale!



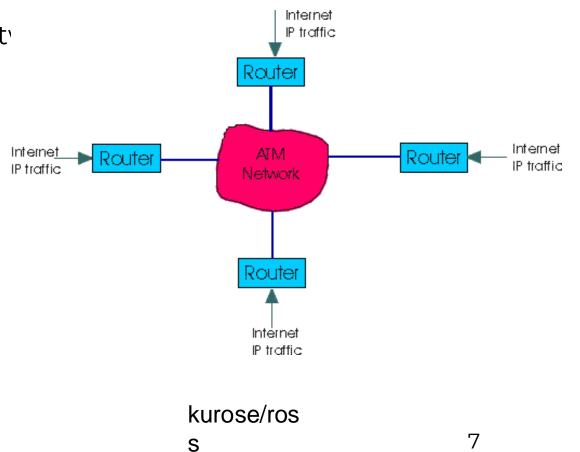
ATM (Carries Cells, not Money)

- ATM (Asynchronous Transfer Mode)
 - Connection-oriented packet-switched network
- Packets are called cells
- 5 byte header + 48 byte payload
- . Fixed length packets are easier to switch in hardware
- . Why?

ATM (Carries Cells, not Money)

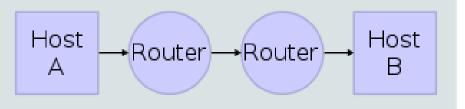
• ATM (Asynchronous Transfer Mode)

- Connection-oriented packet-switched net^v
- Packets are called cells
- 5 byte header + 48 byte payload
- Fixed length packets are easier to switch in hardware
 - Simpler to design
 - Enables parallelism
- . Still used in long distance private links

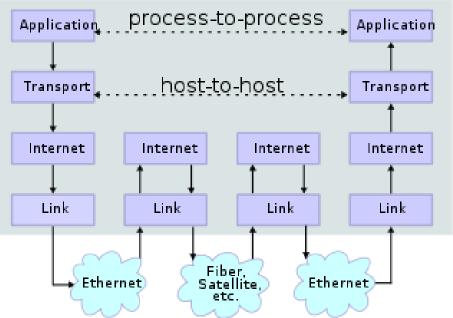


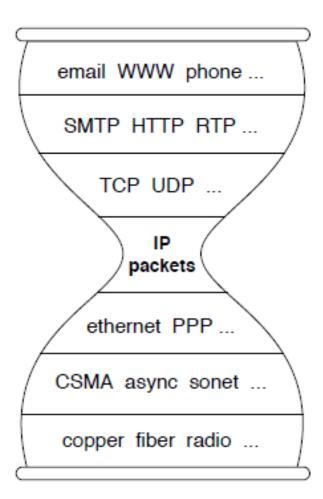
IP Suite – From the First Lecture

Network Topology



Data Flow

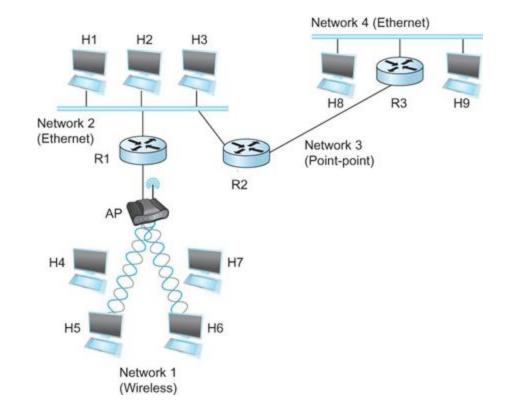




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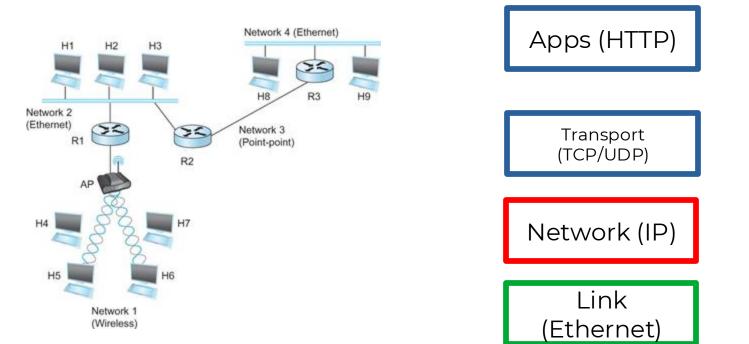
Internet Protocol (IP)

- What is an internetwork?
 - An arbitrary collection of networks interconnected to provide some sort of hosthost to packet delivery service



But that's what switches are for – No?

- Switches create networks, Routers connect different networks.
- Typically switches are at Layer 2, Routers are at Layer 3
- Switches forward FRAMES, Routers forward PACKETS



But that's what switches are for – No?

- This room → Point-to-point link
- This room + next room → Switch
- This room + next room + foundation hall → Switches with VLAN
- This university + Internet \rightarrow Router
- Good for conceptualization not always as simple

Every device has a MAC – Why do we need another address?

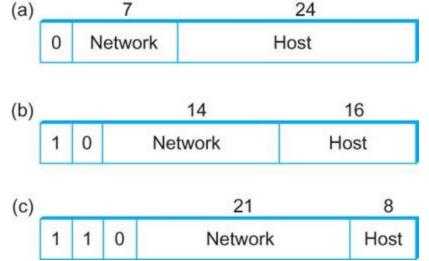
- Ethernet (MAC) addresses are flat
- Not the only link layer
- Not related to network topology
 - Remember we are still connecting to hosts!

. Other reasons?

Apps (HTTP)
Transport (TCP/UDP)
Network (IP Address)
Link (MAC Address)

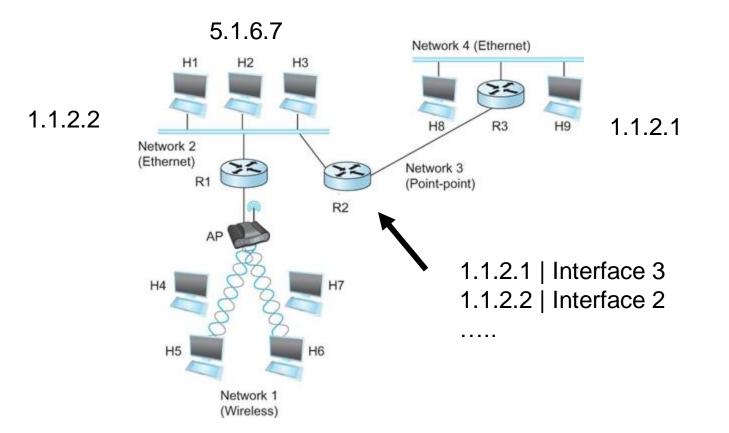
Global Address in IP – Each node has an unique address

- A 32 bit number in quad-dot notation
- . Identifies an Interface
 - . A host might have several interfaces!!!
- 129.82.138.254
 1000001.01010010.10001010.1111110



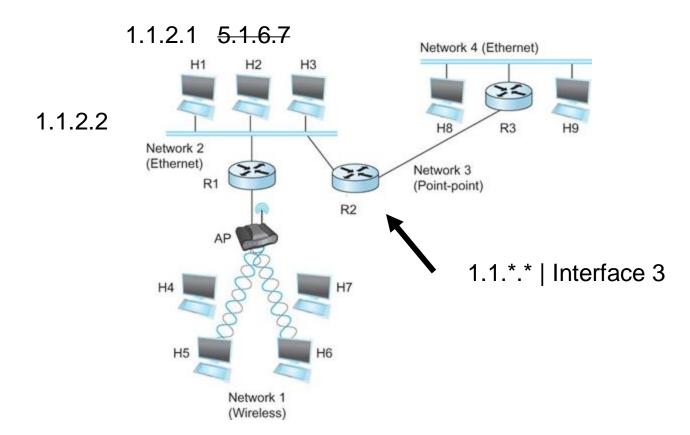
IP allows the network to scale!

• What if addresses were arbitrary?



Solution - Group hosts

• What if addresses were arbitrary?



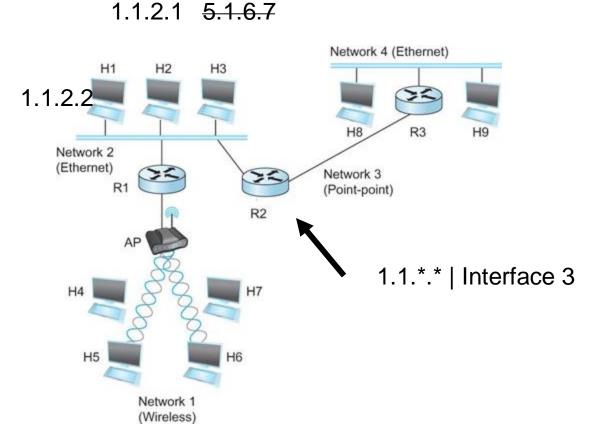
IP addresses are in Network + Host

- 1.1.2.1 →
 - . 1.1 → Network part
 - 2.1 \rightarrow host part
- Each octet can range from 1-255
- Hierarchical address

129.82.138.254

1000001.01010010.10001010.11111110

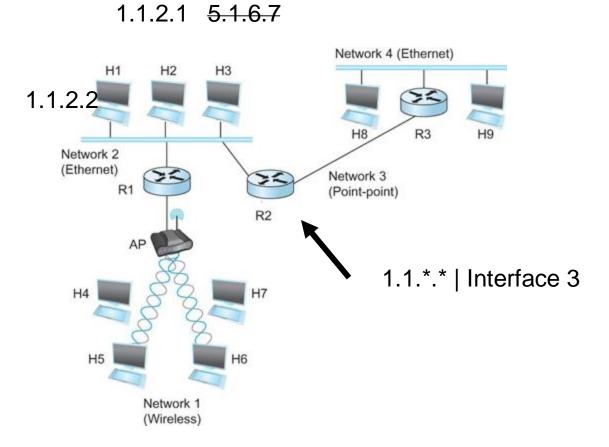
Network part (24 bits). Host part(8 bits)



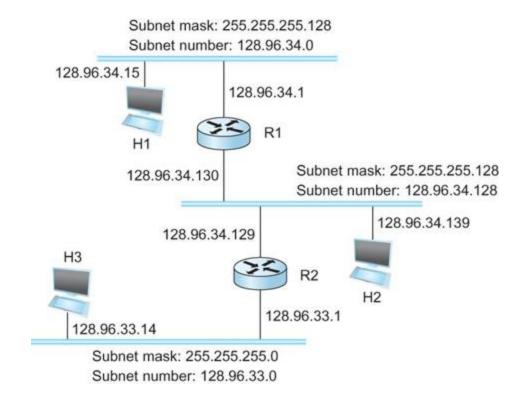
How do we know host vs network → Subnetting

129.82.138.254 (Address)

255.255.255.0 (Subnet mask)



Subnetting

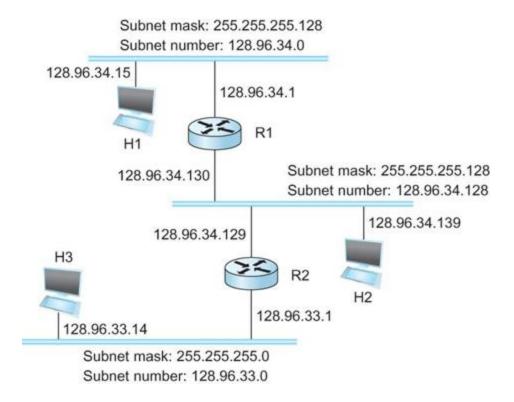


Forwarding Table at Router R1

SubnetNumber	SubnetMask	NextHop
128.96.34.0	255.255.255.128	Interface 0
128.96.34.128	255.255.255.128	Interface 1
128.96.33.0	255.255.255.0	R2

Subnetting

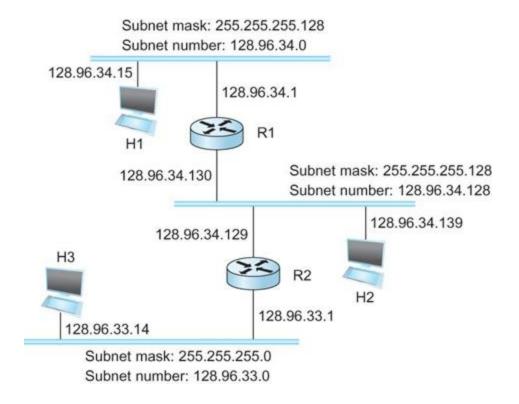
Three classes: Class A: 129.0.0.0/8 Class B: 129.82.0.0/16 Class C: 129.82.2.0/14



SubnetNumber	SubnetMask	NextHop
128.96.34.0	255.255.255.128	Interface 0
128.96.34.128	255.255.255.128	Interface 1
128.96.33.0	255.255.255.0	R2

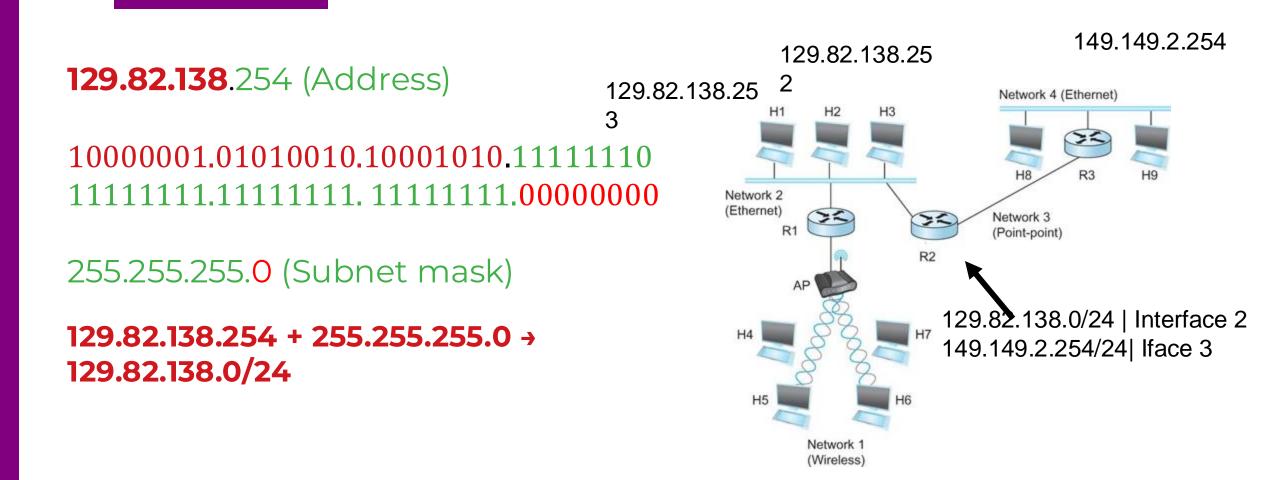
Well, not really!

- CIDR: Classless Interdomain routing
- subnet portion of address of arbitrary length
- address format: a.b.c.d/x, where x is # bits in subnet portion of address
 - 129.82.13.0/23

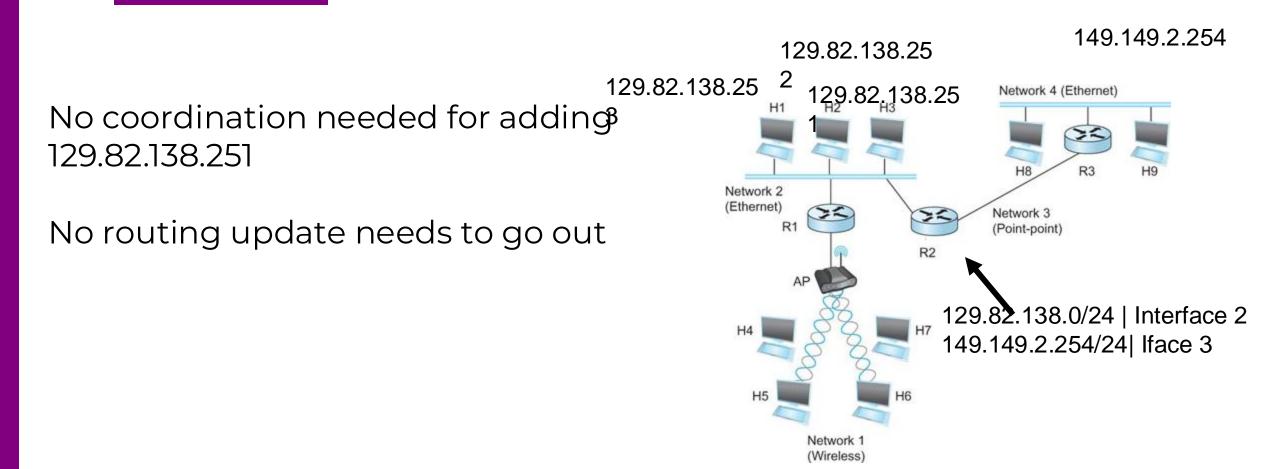


SubnetNumber	SubnetMask	NextHop
128.96.34.0	255.255.255.128	Interface 0
128.96.34.128	255.255.255.128	Interface 1
128.96.33.0	255.255.255.0	R2

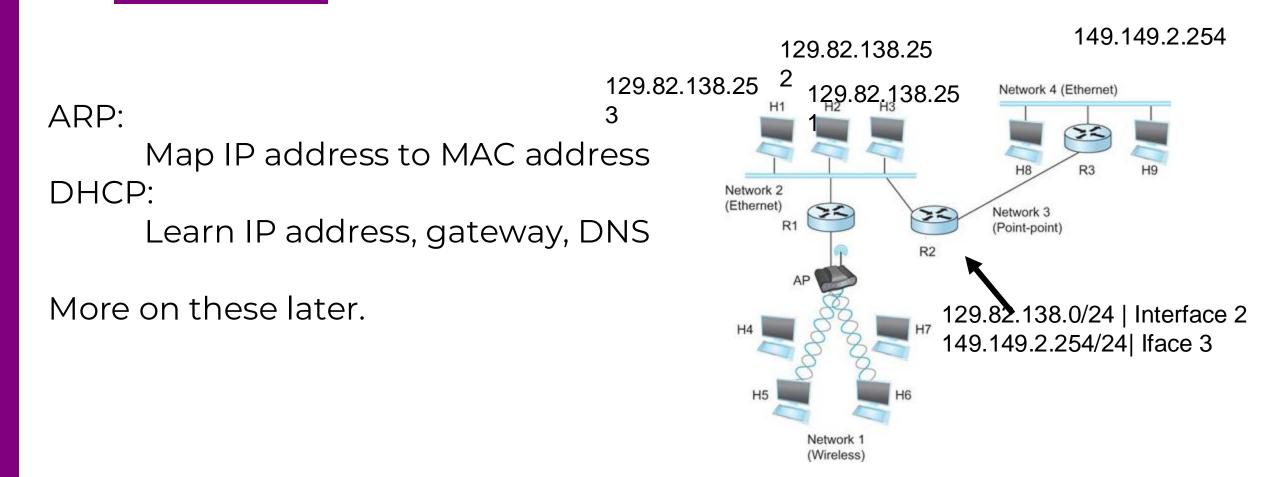
Now routers can operate on Network address!!!!



Address management is localized



Address management can be automated



You have an address – Send data now. IP service model

- Packet Delivery Model
 - Connectionless model for data delivery
- Best-effort delivery (unreliable service)
 - packets are lost
 - packets are delivered out of order
 - . duplicate copies of a packet are delivered
 - packets can be delayed for a long time
- Global Addressing Scheme
 - Provides a way to identify all hosts in the network

IP Packet

Version HLen TOS	Length gs Offset	
Ident Flag	os Offset	
	go	
TTL Protocol	Checksum	
SourceAdd	r	
DestinationAc	ddr	
Options (variable)	Pad (variable)	
Data		

Version (4): 4

Hlen (4): number of 32-bit words in header

TOS (8): type of service (not widely used)

Length (16): number of bytes in this datagram

Ident (16): used by fragmentation

Flags/Offset (16): used by fragmentation

TTL (8): number of hops this datagram has traveled

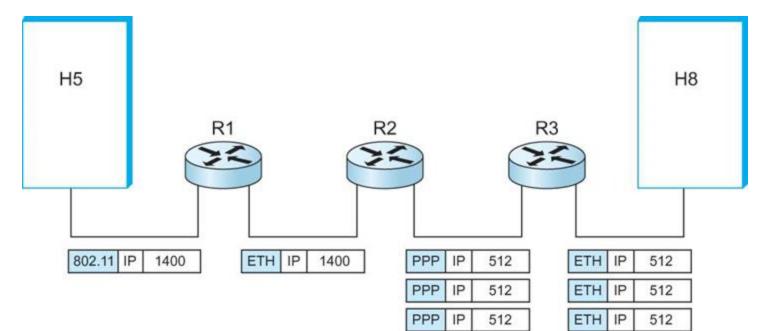
Protocol (8): demux key (TCP=6, UDP=17)

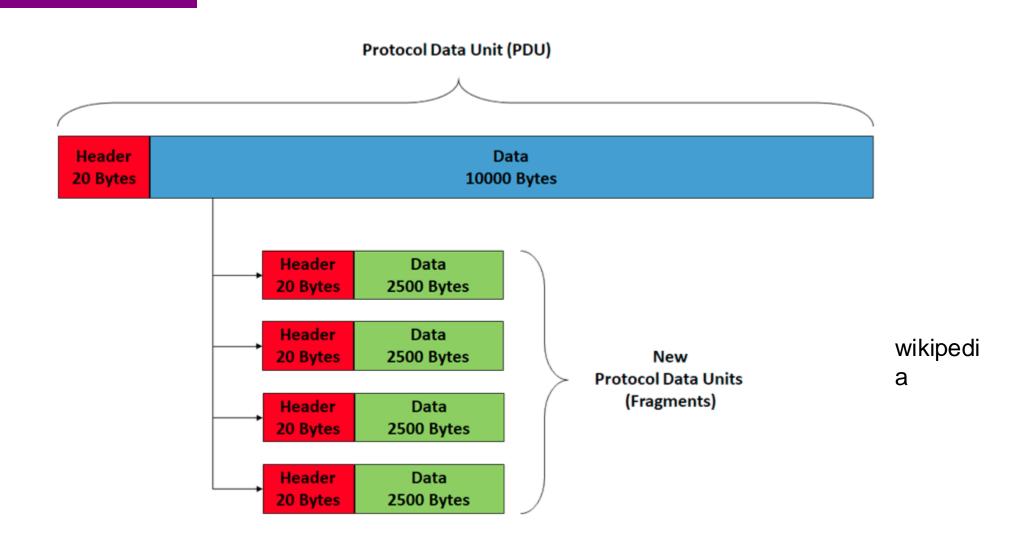
Checksum (16): of the header only

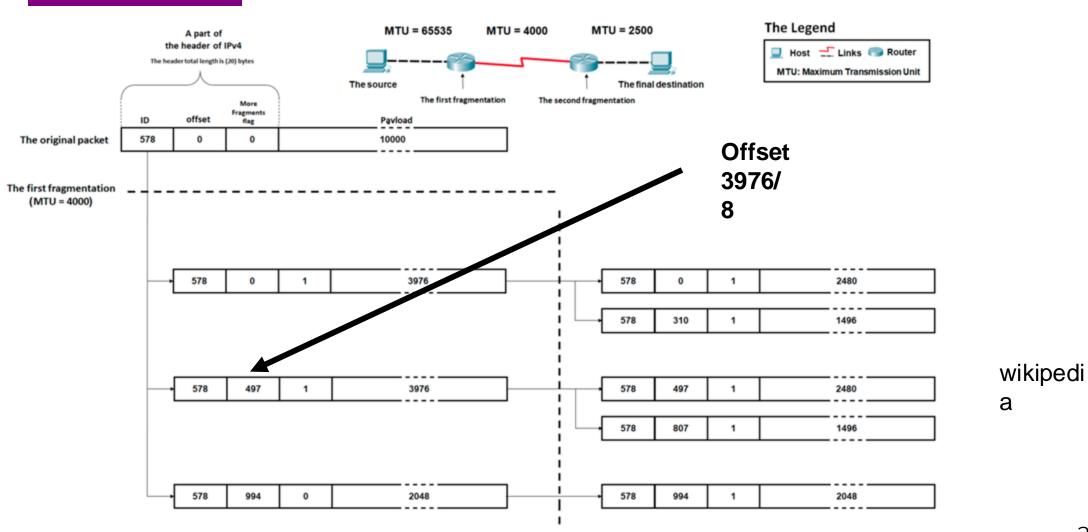
DestAddr & SrcAddr (32)

Underlying Layer 2 limitations

- Ethernet 1500
- PPP 512
- Break packets into smaller chunk and reassemble later



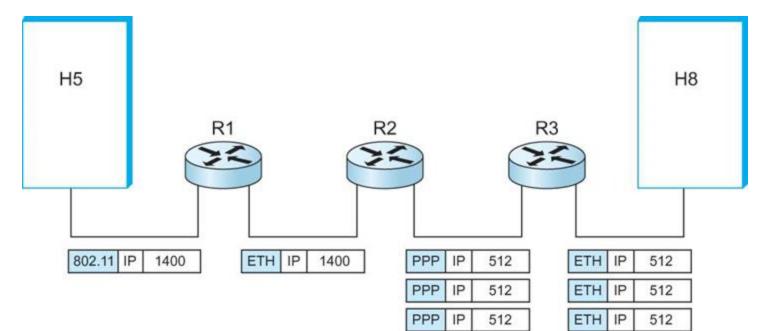




The second fragmentation

Underlying Layer 2 limitations

- Ethernet 1500
- PPP 512
- Break packets into smaller chunk and reassemble later



Reading Assignments

Internetworking: <u>https://book.systemsapproach.org/internetworking/basic-ip.html#what-is-an-internetwork</u>

Upto Global Addresses:

https://book.systemsapproach.org/internetworking/basic-ip.html#global-addresses